SEMISTABLE REDUCTION OVER THICK LOG POINTS

ALEXANDER E. MOTZKIN AND MICHAEL TEMKIN

ABSTRACT. We establish a version of a semistable reduction theorem over a log point with a non-trivial nilpotent structure. In order to do this we extend the classical desingularization theories to non-reduced schemes with generically principal nilradical.

1. INTRODUCTION

The main goal of this work is to establish a version of semistable reduction (or relative desingularization) of schemes over a base of the form $B_n = \operatorname{Spec}(k[\pi]/(\pi^n))$, where k is a field of characteristic zero. In particular, we were asked about existence of such a result by Christian Schnell and Junchao Shentu, and an answer to this problem might have applications to Hodge theory. A general theory of resolution of relative (log) morphisms $X \to B$ was recently developed in [ATW20] under the assumption that B is log regular, so it cannot be directly applied to our case. There is a way to generalize that theory to the base we need, but it will be worked out elsewhere, and in this work we decided to use more standard tools. In fact, we manage to deduce our main result from the classical resolution and principalization theorems.

1.1. Log smoothness over thick points. We would like to stress that a priori it is not even so clear how to formulate an appropriate semistable reduction conjecture whose goal is to resolve a generically smooth morphism $X \to B = B_n$. For comparison, recall the situation with semistable reduction over a trait, say over $S = \operatorname{Spec}(R)$ with $R = k[\pi]_{(\pi)}$. If $Z \to S$ has a connected but not irreducible closed fiber Z_s , for example, $Z = \operatorname{Spec}(R[x,y]/(xy-\pi))$, then by Zariski's connectedness theorem the same will be true for any modification Z' of Z, and hence $Z' \to S$ is not smooth. The best one can do is to resolve Z and make Z_s an snc divisor. In the zero characteristic case this improves the morphism $Z \to S$ drastically: étale-locally Z becomes isomorphic to the model charts of the form $\operatorname{Spec}(R[t_1,\ldots,t_n]/(t_1^{d_1}\ldots t_n^{d_n}-\pi))$, in particular, it is log smooth with respect to the natural closed fiber log structures. For the sake of completeness, we note that such a morphism is called semistable if $d_i \in \{0, 1\}$, and one can further make $Z \to S$ semistable by an extension of R and log (or toroidal) blowings up of Z, but we will not need this.

Adopting the logarithmic approach is also absolutely critical for the results of [ATW20], so it is natural to adopt it in our case too and to provide B with the log structure induced by $\mathbb{N} \log \pi$. Nevertheless, it turns out that one cannot in general modify a generically smooth B-scheme X to X' which is log smooth over B. Here

Key words and phrases. Semistable reduction, thick manifolds.

M.T. was supported by ERC Consolidator Grant 770922 - BirNonArchGeom.

is the simplest example: $X = \operatorname{Spec}(k[x, y]/(y^n))$ with the morphism $X \to B$ given by $\pi = xy$. Indeed, the morphism is not smooth at x = 0 because the nilradical of \mathcal{O}_X is not generated by π at this point. The situation does not improve if we increase the log structure on X by adding $\mathbb{N} \log x$. Moreover, any modification $X' \to X$ induces an isomorphism of the reductions and only increases the nilpotent structure (for example, by replacing y by $y' = y/x^n$), hence it does not improve the situation either. The only way to achieve some sort of resolution of $X \to B$ is to use the whole class of smooth B-curves, including the reducible ones, and there are plenty of them. For example, $Y = \operatorname{Spec}(k[x,y]/(x^ny^n))$ is log smooth over B (where $\pi = xy$), and X is an irreducible component of Y. In fact, this example is the closed fiber of the morphism $Z \to S$ from the example above, but reducible closed fibers over S show up naturally, while reducible schemes over B at first glance look as a pathology, and by no means they can show up as modifications of irreducible ones.

We hope that the above discussion and relation to the classical semistable reduction over S give enough motivation for our definition: a desingularization of a B-scheme X is a modification $X' \to X$ such that locally X' can be realized as an irreducible component and a monomial subscheme of a log smooth B-scheme Y. The main result of the paper – Theorem 3.3.3 proves that any generically smooth B-scheme X possesses such a desingularization. Moreover, we show that it can be chosen of a special combinatorial type: X' is étale over $\operatorname{Spec}(R[\varepsilon, t_1, \ldots, t_n]/(\varepsilon^n))$ with π mapped to $\varepsilon t_1^{d_1}, \ldots, t_n^{d_n}$ and the log structure generated by ε and t_1, \ldots, t_n . In addition, we show that any closed subscheme $Z \hookrightarrow X$ can be made a monomial subscheme of X', which is an analogue of resolution of pairs.

Remark 1.1.1. Here is another argument in favor of our definition of desingularization over B. We mentioned that principalization of ideals on log smooth B-schemes is possible (and will be developed elsewhere). As usually, it can be used to obtain a desingularization of B-schemes X as follows. Separate components, reducing to the case when X is irreducible. Locally embed X into a log smooth B-scheme M. Principalize \mathcal{I}_X on M by blowing up centers which are monomial or log smooth over B and stop when the new blowing up center V contains the generic point of the strict transform X' of X. In particular, V is log smooth over B and X' is its irreducible component. In the situation of [ATW20] (similarly to the classical resolution) the connected components of V are irreducible, hence X' is one of them and $X' \to X$ is a B-resolution. In our situation, a log smooth V does not have to be irreducible, hence we only obtain a modification $X' \to X$ which is an irreducible component of a log smooth V. This is precisely the above definition of a B-desingularization.

Remark 1.1.2. (i) One can wonder if there exist another type of resolutions: a proper morphism $X' \to X$ which is birational on one component of X' and such that $X' \to B$ is log smooth. We do not study this question in the paper, but it seems plausible that in general this is impossible.

(ii) One can also ask if the embedding of X' as an irreducible component of a *B*-log smooth *Y* can be made globally. This also seems to be a not so natural question, though we show in the paper that this is possible, and if $X \to B$ is proper, then *Y* can also be taken to be proper over *B*. This construction uses log blowings up, which are natural non-birational proper morphisms $X' \to X$ (which contract components). However, log blowings up cannot be used to attack the question in (i) because they are log étale, and hence if $X \to B$ is not log smooth, then the same is true for $X' \to B$.

1.2. Outline of the paper. Now, let us briefly describe our other results and the structure of the paper. In Section 2 we extend some resolution results to non-reduced schemes. For the sake of completeness we recall in §2.1 Hironaka's resolution of non-reduced schemes by *thick manifolds* – schemes X with a regular reduction $\tilde{X} = X^{\text{red}}$ and normally flat along \tilde{X} . However, we only need to work with schemes X whose nilradical is generically principal, and their resolution is deduced in §2.3 from the classical functorial principalization: in Theorem 2.3.1 we show that any such X possesses a functorial modification X', which is a principally thick manifold or a ptm – a thick manifold with a locally principal nilradical, see §2.2.

In §3.1 we deduce principalization of subschemes of a ptm X from the classical principalization of subschemes of \widetilde{X} , and in §3.3 we prove our main result on resolution of varieties over B. As in the case of S-schemes, we just resolve X (by a thick manifold since X is non-reduced) and then principalize the divisor (π) on X. There are small additional subtleties in this case, which are dealt with in §3.2. Finally, in §3.4 we explain how analogous results can be proved for certain quasi-excellent schemes, formal schemes and analytic spaces.

Section 4 is devoted to proving that in the case of varieties, after an additional work the resolved *B*-scheme X' can be globally embedded into a log smooth *B*-scheme *Y*. This operation is less functorial, it depends on the choice of a retract, and our argument does not apply to general qe schemes or analytic varieties, because retract does not have to exist in these cases. It might be the case that the generalization of our main result we prove in Theorem 4.4.1 addresses a not so natural question. Still, the intermediate results we prove about retracts and factorization of birational morphisms of non-reduced schemes can be of an independent interest.

2. Desingularization on principally thick manifolds

In this section we will extend the classical resolution theory to schemes which do not have to be generically reduced but whose nilradical is generically principal. Everything will be deduced rather straightforwardly from the classical theory.

2.1. **Hironaka's resolution of non-reduced schemes.** In the famous 1964 paper [Hir64] Hironaka also established desingularization of arbitrary non-reduced varieties that we are going to recall briefly.

2.1.1. *Thick manifolds.* Let us fix some terminology first. We warn the reader that the notion of thick manifolds is non-standard.

Definition 2.1.2. (i) A scheme X is normally flat along a subscheme Y given by an ideal $\mathcal{I} \subseteq \mathcal{O}_X$ if $\mathcal{I}^n/\mathcal{I}^{n+1}$ is a flat module over $\mathcal{O}_Y = \mathcal{O}_X/\mathcal{I}$ for any n.

(ii) A locally noetherian scheme X is called a *thick manifold* if its reduction X^{red} is regular and X is normally flat along X^{red} .

(iii) By thickness of a noetherian scheme X at a point x we mean the sum $h_X(x)$ of dimensions of the k(x)-vector spaces $(\mathcal{N}^n/\mathcal{N}^{n+1}) \otimes k(x)$, where \mathcal{N} is the nilradical of \mathcal{O}_X .

The normal-flatness condition means that although there might be a non-reduced structure, it behaves "as nice as possible".

Lemma 2.1.3. (i) If $\eta \in X$ is a generic point, then $h_X(\eta)$ is the length of the Artin ring $\mathcal{O}_{X\eta}$.

(ii) The thickness function $h_X \colon X \to \mathbb{N}$ is upper-semicontinuous, and X is normally flat along X^{red} if and only if h_X is a locally constant function. In the latter case, X has no embedded components.

Proof. The first claim is obvious. In (ii) set $Y = X^{\text{red}}$ and note that for a finite \mathcal{O}_Y -module M the function $r_M(x) = \dim_{k(x)} M \otimes k(x)$ is upper semicontinuous on Y, and r_M is locally constant if and only if M is flat. The equivalence in (ii) follows by applying this to $h_X = \sum_n r_{\mathcal{N}^n/\mathcal{N}^{n+1}}$ – due to the semicontinuity, h_X is locally constant if and only if each summand is. Finally, if X has an embedded component, then there exists a closed immersion $X' \hookrightarrow X$ such that the locus V where $X' \neq X$ is nowhere dense but non-empty. In this case, $h_X = h_{X'}$ outside of V and h_X jumps on V. In particular, for any generic point η with a specialization $x \in V$ one has $h_X(\eta) = h_{X'}(\eta) \leq h_{X'}(x) < h_X(x)$ and hence X is not a thick manifold.

2.1.4. *Hironaka's theorem.* In a sense, thick manifolds are non-reduced schemes which are as close to regular ones as possible, and desingularization of non-reduced schemes by thick manifolds is the best one can hope for - one has to keep the nilpotent structure at the generic points, but it becomes as simple as possible everywhere else.

Definition 2.1.5. (i) Throughout this paper a morphism $X' \to X$ is called a *modification* if it is proper and induces an isomorphism U' = U of dense open subschemes.

(ii) A desingularization of a scheme X is a modification $X' \to X$ such that X' is a thick manifold.

Hironaka proved that this desingularization is, indeed, possible.

Theorem 2.1.6 ([Hir64, Main Theorem I^{*}]). There exists a construction which associates to any variety X of characteristic zero a projective desingularization $\mathcal{F}(X) = X' \to X$

In fact, Hironaka proved his theorem for the more general class of schemes of finite type over a local quasi-excellent ring of characteristic zero.

2.1.7. Strong resolution. Hironaka's proof is existential, but modern methods that grew from his approach can also provide a smooth-functorial desingularization. Unfortunately, a convenient reference is missing in the literature, so let us describe the situation briefly. Nowadays there are two main resolution methods, both originating from Hironaka's work, that will be called *basic* (or weak in the literature) and *strong*. Basic resolution is achieved by blowing up smooth centers in the ambient manifold M, but the induced blowings up of the scheme X itself may have singular centers. Its primary invariant is the order of an ideal $I_X \subset \mathcal{O}_M$. This resolution was studied in the majority of recent works and it is smooth-functorial by results of Włodarczyk, see [Wło05].

Strong resolution uses the whole Hilbert-Samuel function as the primary invariant and it only blows up smooth centers such that X is normally flat along them.

The price one has to pay is that the process is much heavier, as well as the proofs. In fact, one reduces this resolution to the basic one by encoding the HS function in terms of the order of another ideal (called a presentation). So far, the only known way to resolve general non-reduced schemes is via a strong resolution methods. A canonical strong desingularization was established by Bierstone-Milman, see [BM97, Theorem 11.14] (though it was only formulated in the embedded version), and its functoriality properties were briefly checked in [BMT11, Theorem 6.1. Addendum] (one should check that the reduction to a basic method is smooth-functorial).

However, we will see that already the basic method suffices to resolve varieties with generically principal nilradical, and this is the only case we will need later. In particular, smooth-functoriality is then automatic.

2.2. Principally thick manifolds. We say that a scheme X is *nilprincipal* if its nilradical is a locally principal ideal. If, in addition, X is a thick manifold, then we say that X is a *principally thick manifold* or a *ptm*.

Definition 2.2.1. Let X be a ptm. By a family of parameters at a point $x \in X$ we mean a family $\varepsilon, t_1, \ldots, t_d \in \mathcal{O}_{X,x}$ such that ε generates the nilradical of $\mathcal{O}_{X,x}$ and t_1, \ldots, t_d induce a family of regular parameters of $\mathcal{O}_{X^{\text{red}},x} = \mathcal{O}_{X,x}/(\varepsilon)$. We call ε a nilpotent parameter and call t_1, \ldots, t_d regular parameters.

Thickness on a ptm has especially simple interpretation.

Lemma 2.2.2. Let X be a ptm, $x \in X$ a point and $\varepsilon \in \mathcal{O}_x$ a nilpotent parameter at x. Then the thickness $h_X(x)$ is the minimal number h such that $\varepsilon^h = 0$.

Proof. Let \mathcal{N} be the nilradical, then $\mathcal{N}_x = \varepsilon \mathcal{O}_x$ and $\mathcal{N}^n / \mathcal{N}^{n+1}$ is non-zero if and only if $\mathcal{N}^n \neq 0$, that is n < h. In the latter case, $(\mathcal{N}^n / \mathcal{N}^{n+1}) \otimes k(x)$ is generated by the image of ε^n and hence is of dimension 1.

Remark 2.2.3. (i) Due to our definition, ptm's of thickness 1 are just usual manifolds (or regular schemes) with dummy nilpotent parameter $\varepsilon = 0$. Because of this they should often be dealt with slightly differently, and since all results for them are standard we will simply exclude this case when this is convenient.

(ii) We will later enhance ptm's to log schemes with the log structure generated by $\log(\varepsilon)$. In this situation, the case of thickness 1 is not exceptional anymore, in particular, the corresponding log scheme is not log regular.

Lemma 2.2.4. If X is a non-reduced ptm and $\varepsilon, t_1, \ldots, t_d$ is a family of parameters at x, then their images form a basis of the cotangent space m_x/m_x^2 .

Proof. Since X is singular and of dimension d at x, we have $\dim(m_x/m_x^2) > d$. The d+1 parameters generate m_x , hence their images generate m_x/m_x^2 and by a dimension consideration they form a basis.

Lemma 2.2.5. Let M be a regular scheme and $X \hookrightarrow M$ a closed subscheme given by an ideal $\mathcal{I} \subset \mathcal{O}_M$, $x \in X$, $\dim_x(M) = d$ and $\dim_x(X) = l$. Then X is a non-reduced ptm locally at a point $x \in X$ if and only if there exist a family of local parameters $x_1, \ldots, x_d \in \mathcal{O}_{M,x}$ such that $\mathcal{I}_x = (x_1^n, x_2, \ldots, x_r)$ for r = d - l. In addition, if these conditions are satisfied and $(\varepsilon, t_1, \ldots, t_l)$ is a family of parameters of $\mathcal{O}_{X,x}$, then one can choose $x_1, x_{r+1}, \ldots, x_d$ to be any lift of $(\varepsilon, t_1, \ldots, t_l)$.

Proof. The inverse implication is obvious since X^{red} is locally given by the vanishing of (x_1, \ldots, x_r) , and the nilradical of $\mathcal{O}_{X,x}$ is generated by the image of x_1 .

Conversely, assume that X is a ptm and fix any family $(x_1, x_{r+1}, \ldots, x_d)$ of elements of $\mathcal{O}_{M,x}$ that reduces to a family of parameters $(\varepsilon, t_1, \ldots, t_l)$ at x. The images of $(x_1, x_{r+1}, \ldots, x_d)$ in $m_{M,x}/m_{M,x}^2$ are linearly independent because their images in $m_{X,x}/m_{X,x}^2$ form a basis by Lemma 2.2.4. The kernel of the map $m_{M,x}/m_{M,x}^2 \twoheadrightarrow m_{X,x}/m_{X,x}^2$ is the image of \mathcal{I}_x , hence the family $(x_1, x_{r+1}, \ldots, x_d)$ can be completed to a family of regular parameters (x_1, \ldots, x_d) so that $x_2, \ldots, x_r \in \mathcal{I}_x$.

Let *n* be the thickness of *X* at *x*. Then locally at *x* the ideal $\mathcal{I}' = (x_1^n, x_2, \ldots, x_r)$ is contained in \mathcal{I} , hence defines a ptm *X'* containing *X*. We claim that the surjection $\phi \colon \mathcal{O}_{X',x} \twoheadrightarrow \mathcal{O}_{X,x}$ is an isomorphism. Indeed, any element $a \in \text{Ker}(\phi)$ is nilpotent because X^{red} and X'^{red} are smooth of the same dimension l = d - r at *x*, and using that *X* and *X'* are of thickness *n* we obtain that a = 0.

We say that a ptm (or a scheme) X is *locally embeddable (in a regular scheme)* if any point $x \in X$ has a neighborhood isomorphic to a closed subscheme of a regular scheme M. For example, any k-variety satisfies this property. The above lemma implies that locally embeddable ptms are isomorphic to a divisor nD with a regular D in an ambient regular scheme:

Corollary 2.2.6. If X is a locally embeddable ptm, then any point $x \in X$ possesses a neighborhood U and a closed immersion $U \hookrightarrow M$, such that M is regular and U is the divisor in M of the form $V(t^n)$, where t is a parameter of $\mathcal{O}_{M,x}$ and n is the thickness of X at x.

Proof. By Lemma 2.2.5, x has a neighborhood U isomorphic to a closed subscheme $V(t_1^n, t_2, \ldots, t_r)$ of a regular scheme M', where t_1, \ldots, t_d is a regular family of parameters of $\mathcal{O}_{M',x}$. It remains to take $M = V(t_2, \ldots, t_r)$ and define $t \in \mathcal{O}_{M,x}$ to be the image of t_1 .

The following corollary certainly holds for arbitrary ptm's, but we only need the case of locally embeddable ones and then the argument is very simple.

Corollary 2.2.7. Any ptm X, which is locally embeddable in a regular scheme, is Cohen-Macaulay. In particular, if $i: U \hookrightarrow X$ is an open immersion, whose complement is of codimension at least 2, then $\mathcal{O}_X = i_*\mathcal{O}_U$.

Proof. By Corollary 2.2.6, locally X is a Cartier divisor in a smooth (hence CM) variety. Therefore, X is itself CM. The second claim is, perhaps, better known for normal schemes, but it holds more generally for S_2 -schemes and, in fact, provides an equivalent characterization of this property. The original source is the theory of Z-purity, see [Gro67, Théorème IV.5.10.5], a modern treatment (in the relative setting) is in [HK04, Proposition 3.5].

2.2.8. *Other characterizations*. For the sake of completeness we provide alternative characterizations of ptm's, though this will not be used in the sequel.

Lemma 2.2.9. Given a scheme X the following conditions are equivalent:

- (i) X is a ptm.
- (ii) X is a generically nilprincipal thick manifold.
- (iii) X^{red} is regular, X is nilprincipal and without embedded components.

Proof. (i) implies (ii) and (iii) because thick manifolds have no embedded components by Lemma 2.1.3(ii). In the sequel, let \mathcal{I} denote the nilradical of \mathcal{O}_X .

(ii) \Longrightarrow (i) Looking at the generic points we see that the locally free module $\mathcal{I}/\mathcal{I}^2$ is of rank 1. Its local generator at a point x lifts to a local generator of the stalk \mathcal{I}_x and hence X is nilprincipal.

(iii) \Longrightarrow (i) We will work locally at $x \in X$. Let $\varepsilon \in \mathcal{O}_x$ be a generator of \mathcal{I}_x and let n be the minimal number with $\varepsilon^n = 0$. We should prove that X is normally flat along X^{red} at x and we claim that, in fact, each $\mathcal{I}_x^i/\mathcal{I}_x^{i+1}$ with i < n is free of rank 1 over $\mathcal{O}_{X^{\text{red}},x} = \mathcal{O}_{X,x}/\mathcal{I}_x$. Clearly, $\mathcal{I}_x^i/\mathcal{I}_x^{i+1}$ is generated by the image of ε^i , so we should prove that its annihilator in $\mathcal{O}_{X,x}/\mathcal{I}_x$ is zero. Moreover, it suffices to prove the latter only for i = n - 1, and since $\mathcal{I}_x^n = 0$ this happens if and only if the annihilator of ε^{n-1} in $\mathcal{O}_{X,x}$ is contained in (in fact, equal to) \mathcal{I}_x . The latter holds because otherwise X would contain an embedded component on whose complement ε^{n-1} vanishes.

Remark 2.2.10. Note that X is generically nilprincipal if and only if for any generic point $\eta \in X$ the Artin ring \mathcal{O}_{η} has principal maximal ideal $m_{\eta} = (\varepsilon)$. Such Artin rings are very special, e.g. see [AM69, Proposition 8.8], in addition, they can be described as quotients of DVR's by non-zero ideals.

2.3. **Desingularization of generically nilprincipal varieties.** Now we are ready to formulate the resolution result for non-reduced schemes with simplest nilpotent structure.

Theorem 2.3.1. There exists a construction which associates to any generically nilprincipal variety X of characteristic zero a projective desingularization $\mathcal{F}(X) = X' \to X$ which depends on X smooth-functorially: for any smooth morphism $Y \to X$ one has that $\mathcal{F}(Y) = \mathcal{F}(X) \times_X Y$.

Proof. Step 1. Assume that X is irreducible and embedded as a closed subscheme in a smooth variety M. In this case we apply the smooth-functorial principalization from [Wło05] to $\mathcal{I}_X \subset \mathcal{O}_M$ (one can use also the methods of [Kol07] or [BM08], but they all are equivalent). It outputs a sequence of blowings up $f_i: M_i \to M_{i-1}$, $1 \leq i \leq n$ with smooth centers $V_i \subset M_i$ and snc boundaries $E_i \subset M_i$ such that $E_0 = \emptyset, M_0 = M, E_i = f_i^{-1}(E_{i-1}) \cup f_i^{-1}(V_{i-1}), V_i$ has simple normal crossings with E_i and lies in the preimage of X, and $\mathcal{I}_X \mathcal{O}_{M_n}$ is invertible and supported on E_n . As one always does when deducing resolution from principalization, we cut this sequence at M_l such that f_{l+1} is the first blowing up whose center V_l maps onto X; it exists because the preimage of X lies in E_n . Then the strict transform $X_l \subset M_l$ of X (i.e. the schematic closure of $X \times_M M_l \setminus E_l$) is non-empty, lies in V_l and contains an irreducible component of V_l . Since V_l is smooth, X_l^{red} is an irreducible component of V_l , in particular, it is smooth. So far, we repeated the standard argument, which constructs the resolution $X_l^{\text{red}} \to X^{\text{red}}$ of the integral scheme X^{red} , but we claim that, moreover, $X_l \to X$ is a resolution of X.

Generically on X the principalization works as follows: there is no boundary and whenever $\dim(M) > \dim(X) + 1$ the order is one by Lemma 2.2.5, so one simply replaces M by a maximal contact – any smooth divisor containing X. Once $\dim(M) = \dim(X) + 1$, one has $\mathcal{I}_X = (t^n)$, where n is the thickness of X and H =V(t) is the reduction of X. It follows that the order is n, the only maximal contact is H and the algorithm simply resolves the marked ideal (\mathcal{I}_X, n) by blowing up H. By our assumption this happens at stage l, so at this stage at any point $x \in X_l$ there is no new boundary (we are in Step 1b of the algorithm in [Who05]), $d = \operatorname{ord}_x(\mathcal{I}_{X_l})$ equals 1 whenever $\dim(M_l) > \dim(X_l) + 1$, and d = n once $\dim(M_l) = \dim(X_l)$ because the algorithm first deals with the locus where $\dim(X_l)$ is maximal and by our assumption it blows up the whole $H = X_l^{\text{red}}$. Note that nH is the schematic closure of the generic point of X_l , hence $nH \hookrightarrow X_l$ and \mathcal{I}_{X_l} is contained in the divisorial ideal \mathcal{I}_H^n . In particular, $\mathcal{I}_H^{-n}\mathcal{I}_{X_l}$ is an ideal, and since the degrees of \mathcal{I}_H^n and \mathcal{I}_{X_l} are equal to *n* everywhere along X_l , the order of $\mathcal{I}_H^{-n}\mathcal{I}_{X_l}$ is zero, that is, $X_l = nH$. So, as claimed X_l is a principally thick manifold and $X_l \to X$ is a projective desingularization.

For the sake of completeness, we note that the blowing up $X_{l+1} \to X_l$ is the last blowing up of the principalization, it happens at Step 1ba in [Wh005], and the invariant at this step is $(1, 0, 1, 0, \dots, 1, 0, n, 0, \infty)$.

Step 2. Assume that X is irreducible. In this case one locally embeds X into a smooth variety M and takes the resolution $X_l \to X$ provided by Step 1. A standard argument used in all works on functorial resolution shows that this is independent of the local choice of embeddings and globalizes to a smooth-functorial method to resolve X – usually one works with an integral X because principalization does not provide resolution of general non-reduced varieties, but this reduction only uses that X is irreducible.

Step 3. The general case. In general, let η_1, \ldots, η_s be the generic points of X with the induced scheme structure and let X_i be the schematic closure of η_i . Then $\coprod_i X_i \to X$ is a modification, which depends on X in a way which is easily seen to be compatible with smooth morphisms $X' \to X$. If $Y_i \to X_i$ are the resolutions provided by Step 2, then $\mathcal{F}(X) = \coprod_i Y_i$ is a resolution of X, and by the above \mathcal{F} is smooth-functorial.

3. Resolution on PTM's

In this section we will prove our main result – a version of semistable reduction over closed subschemes of a trait (i.e. zero dimensional schemes with principal radical). In the sequel, given a ptm X we will use the notation $\tilde{X} = X^{\text{red}}$ to denote its reduction.

3.1. **Principalization on ptm's.** In this section we construct principalization of ideals on a ptm using the usual principalization on its reduction.

3.1.1. *Blowings up.* Principalization is achieved by iteratively blowing up regular centers.

Lemma 3.1.2. Assume that X is a ptm and $V \hookrightarrow X$ a regular center. Then $X' = \operatorname{Bl}_V(X)$ is a ptm too, and the reduction of X' is the blowing up of the reduction of X along V, that is, $\widetilde{X}' = \operatorname{Bl}_V(\widetilde{X})$.

Proof. The second claim follows from the fact that \widetilde{X}' is the strict transform of the closed subscheme \widetilde{X} of X, and hence by the classical properties of blowings up \widetilde{X}' is the blowing up of \widetilde{X} along $V \times_X \widetilde{X} = V$. In particular, \widetilde{X}' is regular.

It remains to check that X' is a ptm, and this is a local claim over a point $x \in V$. So, we can assume that X = Spec(A) is a local scheme with closed point x. Choose local parameters $\varepsilon, t_1, \ldots, t_n$ at x, such that $V = V(\varepsilon, t_1, \ldots, t_r)$. Then the ε -chart is empty, so X' is covered by the charts $X'_i = \text{Spec}(A_i)$, with $A_i = A[\varepsilon/t_i, t_1/t_i, \ldots, t_r/t_i] \subseteq A_{t_i}$, and it is easy to see that $\varepsilon' = \varepsilon/t_i$ generates the nilradical of A_i . Thus, the nilradical of X' is locally principal and the reduction \widetilde{X}' is regular, that is, X' is a ptm.

3.1.3. Snc divisors. Working on a ptm X one has to distinguish Weil and Cartier divisors. Weil divisors are just divisors on the reduction, while Cartier divisors contain infinitesimal information as well. Saying a "divisor" on X we will always mean an effective Cartier divisor $D = V(f) \subset X$. A divisor D will be called *snc* if locally on X it factors as $D = \sum_i D_i$ such that each reduction $\widetilde{D}_i = D_i \times_X \widetilde{X}$ is smooth and $\widetilde{D} = \bigcup_i \widetilde{D}_i$ is an snc divisor on \widetilde{X} . A divisor D' is D-monomial if it is of the form $D' = \sum_i n_i D_i$. A regular closed subscheme $V \hookrightarrow X$ has simple normal crossings with \widetilde{D} as a subscheme of \widetilde{X} .

Remark 3.1.4. If D is a divisor with an snc reduction $\widetilde{D} = \bigcup_i D_i$, then it can freely happen that D is not snc because the factorization of \widetilde{D} , does not lift to a factorization of D, and the same is true for monomiality. For example, if X = $\operatorname{Spec}(k[x, y, \varepsilon]/(\varepsilon^2))$ then $D = V(xy + \varepsilon)$ is not snc, though \widetilde{D} is snc, and D' = $V(y^2 + \varepsilon)$ is not V(y)-monomial, though \widetilde{D}' is V(y)-monomial. In addition, there exists different snc divisors with the same reduction, for example, (x) and $(x + \varepsilon)$.

In the sequel, by a *ptm with a boundary* we mean a ptm X and an snc divisor $E \hookrightarrow X$, also called a *boundary*. A boundary E is called *ordered* if it is provided with a decomposition $E = \sum_{i=1}^{l} E_i$ into components (not necessarily irreducible) with smooth reductions \tilde{E}_i .

3.1.5. Addmissible blowings up and transforms. Assume that (X, E) is a ptm with a boundary and $\mathcal{I} \subseteq \mathcal{O}_X$ is an ideal corresponding to a closed subscheme $Z = V(\mathcal{I})$, then a Z-admissible blowing up is a blowing up $f: X' = \operatorname{Bl}_V(X) \to X$, where V is regular, has snc with E and is contained in Z. Since $\mathcal{I} \subseteq \mathcal{I}_V$ and the exceptional divisor $E_f = V \times_X X'$ is a Cartier divisor, the principal transform $\mathcal{I}' = (\mathcal{I}\mathcal{O}_{X'})\mathcal{I}_{E_f}^{-1}$ is defined. It corresponds to the subscheme Z' of X' that we also denote $Z \times_X X' - E_f$ and call the principal transform of Z under f.

In addition, if $E = \sum_{i=1}^{n} E_i$ is ordered, then for each *i* we set $E'_i = E_i \times_X X' - V_i \times_X X'$, where V_i is the union of all components of *V* contained in E_i (and so $V_i \times_X X'$ is the component of E_f contained in $E_i \times_X X'$). Finally, we set $E'_{n+1} = E_f$ and define the total transform of *E* to be the ordered boundary $E' = \sum_{i=1}^{n+1} E'_i$. Since E' is independent of the ordering of *E* and an ordering can be chosen locally, we also obtain a definition of total transform for unordered boundaries.

When the thickness is 1 these definitions are nothing but the classical ones from the principalization theory on regular schemes. Moreover, they are compatible with the reduction:

Lemma 3.1.6. Assume that (X, E) is a ptm with an (ordered) boundary, $Z \hookrightarrow X$ is a closed subscheme and $f: X' \to X$ is a Z-admissible blowing up with center $V \hookrightarrow Z$, and let Z' and E' be the transforms. Also, let $\widetilde{Z} = Z \times_X \widetilde{X}$ and $\widetilde{E} = E \times_X \widetilde{X}$, and let \widetilde{Z}' and \widetilde{E}' be their transforms under the blowing up of \widetilde{X} along V. Then, E' is an snc divisor in X', $\widetilde{Z}' = \widetilde{Z} \times_{X'} \widetilde{X}'$ and $\widetilde{E}' = E' \times_{X'} \widetilde{X}'$.

Proof. Recall that by Lemma 3.1.2 $\tilde{f}: \tilde{X}' \to \tilde{X}$ is the blowing up along V. In particular, $E_f = V \times_X X'$ and $E_{\tilde{f}} = V \times_{\tilde{X}} \tilde{X}'$ and it follows that $E_{\tilde{f}} = V \times_X \tilde{X}' = E_f \times_{X'} \tilde{X}'$. Using that $(Z \times_X X') \times_{X'} \tilde{X}' = Z \times_X \tilde{X}' = \tilde{Z} \times_{\tilde{X}} \tilde{X}'$, we obtain that also $\tilde{Z}' = \tilde{Z} \times_{X'} \tilde{X}'$.

For concreteness assume that $E = \sum_{i=1}^{n} E_i$ is ordered and let $E' = \sum_{i=1}^{n+1} E'_i$ and $\tilde{E}' = \sum_{i=1}^{n+1} \tilde{E}'_i$. The above paragraph easily implies that $\tilde{E}'_i = \tilde{E}_i \times_{X'} \tilde{X}'$ for any *i*. Since by the usual theory \tilde{E}' is an snc divisor with irreducible components \tilde{E}'_i , it follows that E' is snc.

3.1.7. Principalization. Assume that (X, E) is a ptm with a boundary and $Z \hookrightarrow X$ is a closed subscheme, then a principalization of Z (or of its associated ideal \mathcal{I}_Z) on (X, E) is a sequence of Z_i -admissible blowings up $X_{i+1} = \operatorname{Bl}_{X_i}(V_i) \xrightarrow{f_i} X_i$ with $0 \leq i \leq n-1$, where $(X_0, E_0, Z_0) = (X, E, Z)$ and for $0 \leq i < n$ one defines inductively E_{i+1} and Z_{i+1} to be the transforms of E_i and Z_i , and $Z_n = \emptyset$. We will use dashed arrows to denote sequences of morphisms, e.g. $X_n \dashrightarrow X$.

3.1.8. The principalization theorem. It turns out that the usual principalization on \widetilde{X} induces a principalization on X yielding the following result:

Theorem 3.1.9. There exists a method which to any triple (X, E, Z) consisting of a ptm X of finite type over a field k of characteristic zero, a boundary E and a closed subscheme $Z \hookrightarrow X$ associates a principalization

$$\mathcal{P}(X, E, Z) \colon X_n \dashrightarrow X_0 = X$$

so that \mathcal{P} depends on the triple (X, E, Z) smooth-functorially: if $X' \to X$ is a smooth morphism of ptm's, $E' = E \times_X X'$ and $Z' = Z \times_X X'$, then $\mathcal{P}(X', E', Z')$ is obtained by pulling back the blowings up sequence $\mathcal{P}(X, E, Z)$ and removing all trivial blowings up (those with empty centers).

Proof. Choose any smooth-functorial principalization method for varieties, e.g. the method from [Wh005]. Set $\tilde{E} = E \times_X \tilde{X}$ and $\tilde{Z} = Z \times_X \tilde{X}$ and let $\tilde{X}_{i+1} = \text{Bl}_{V_i}(\tilde{X}_i)$, $0 \leq i < n$ be the principalization of \tilde{Z} on (\tilde{X}, \tilde{E}) . We pushforward the blowing up sequence $\tilde{X}_n \dashrightarrow \tilde{X}_0 = \tilde{X}$ by setting $X_0 = X$ and $X_{i+1} = \text{Bl}_{V_i}(X_i)$. This makes sense because applying Lemma 3.1.2 inductively we obtain that \tilde{X}_i is, indeed, the reduction of X_i and hence $V_i \hookrightarrow \tilde{X}_i \hookrightarrow X_i$. Applying Lemma 3.1.6 inductively we obtain that V_i has snc with the boundary E_i , which is the transform of E_{i-1} , and V_i is a subscheme of Z_i , which is the transform of Z_{i-1} . In other words, the blowings up are Z_i -admissible. Finally, Lemma 3.1.6 also implies that $\tilde{Z}_i = Z_i \times_{X_i} \tilde{X}_i$ for any $i \leq n$. Since \tilde{Z}_n is empty this implies that Z_n is empty, and hence the sequence $X_n \dashrightarrow X_0 = X$ is a principalization of Z, that we denote $\mathcal{F}(X, E, Z)$.

Smooth-functoriality of the constructed principalization follows from the smooth-functoriality of the usual principalization and the easily checked fact that all other ingredients, such as the reduction morphism $\widetilde{X} \to X$, are compatible with smooth base changes.

3.2. Divisors with monomial reduction.

3.2.1. Modifications with a trivial reduction. Let $f: Y \to X$ be a morphism of ptm's. Then the following conditions are equivalent: (i) f is a bijective modification, (ii) is a modification and a homeomorphism, (iii) f is generically an isomorphism and the reduction \tilde{f} is an isomorphism. If these conditions are satisfied, we say that f is a modification with a trivial reduction.

Example 3.2.2. An important example of a modification with a trivial reduction is the blowing up $Y = \operatorname{Bl}_{\widetilde{D}}(X)$ along a divisor $\widetilde{D} \hookrightarrow \widetilde{X}$ of the reduction. Locally on X we can assume that $X = \operatorname{Spec}(A)$ is affine and $f \in A$ is such that its image $\widetilde{f} \in A/\mathcal{N}$ defines \widetilde{D} , where \mathcal{N} is the nilradical. Then $\widetilde{D} = V(\mathcal{N} + (f))$ in X and $Y = \operatorname{Spec}(A[\frac{N}{f}])$ consists of the f-chart, since other charts correspond to nilpotent elements and hence are empty.

More generally, if $\widetilde{D}_1, \ldots, \widetilde{D}_l$ are divisors of \widetilde{X} then by induction on l one can define a sequence $X_l \dashrightarrow X_0 = X$ with $X_{i+1} = \operatorname{Bl}_{\widetilde{D}_i}(X_i)$ because $\widetilde{D}_i \subset \widetilde{X} = \widetilde{X}_i$.

3.2.3. Monomialization. We saw in Remark 3.1.4 that if D is a divisor with an \tilde{E} -monomial reduction, D itself does not have to be monomial with respect to any boundary or it can be monomial with respect to another boundary. All these problems are remedied by applying a simple blowing up sequence.

Lemma 3.2.4. Assume that X is a ptm with an ordered boundary E and D is a Cartier divisor such that \widetilde{D} is \widetilde{E} -monomial. Let $\widetilde{D} = \sum_{i,n} n\widetilde{D}_{i,n}$ be the decomposition, where $\widetilde{D}_{i,n}$ is the union of components of D of multiplicity n that lie in \widetilde{E}_i . Let $X' \to X$ be the modification with trivial reduction obtained by blowing up centers $V_{1,1}, \ldots, V_{1,n_1}, V_{2,1}, \ldots$, where $V_{i,j} = \prod_{n \geq j} \widetilde{D}_{i,n}$. Then $D \times_X X'$ is monomial with respect to the transform E' of E.

Proof. It suffices to check the claim locally at a point $x \in X$. Let $n_i \geq 0$ be such that $x \in \widetilde{D}_{i,n_i}$. Choose parameters $\varepsilon, t_1, \ldots, t_l$ so that $E = V(t_1 \ldots t_r)$ at x. Then locally at x we have that $\widetilde{D} = V(\prod_{i=1}^r \widetilde{t}_i^{n_i})$ and hence D = V(f) for $f = \prod_{i=1}^r t_i^{n_i} + \varepsilon a$. Since $\varepsilon' = \varepsilon / \prod_{i=1}^r t_i^{n_i}, t_1, \ldots, t_n$ are parameters of X' at x we obtain that $f = \prod_{i=1}^r t_i^{n_i} (1 + \varepsilon' a)$ in $\mathcal{O}_{X',x}$ and $1 + \varepsilon' a$ is a unit. This finishes the proof.

For the sake of completeness we note that the same sequence is produced by the usual principalization:

Remark 3.2.5. In fact, the sequence $\widetilde{X}' \dashrightarrow \widetilde{X}$ in the above lemma is precisely the principalization blowings up sequence of \widetilde{D} on $(\widetilde{X}, \widetilde{E})$ produced by the classical order reduction algorithms with marking 1. So, similarly to the argument in the proof of Theorem 3.1.9, its pushforward $X' \dashrightarrow X$ in the lemma is a principalization of D.

3.3. Resolution of *B*-schemes.

3.3.1. Distinguished B-pairs. Let $B = \operatorname{Spec}(k[\pi]/(\pi^n))$. By a distinguished Bpair we mean a morphism $f: X \to B$ and a divisor E on X such that X is a ptm, E is a boundary on X, and $\pi \mathcal{O}_X = \mathcal{NI}$, where \mathcal{N} is the nilradical of \mathcal{O}_X and $V(\mathcal{I})$ is an E-monomial divisor. In down to earth terms this means that locally on X there exists coordinates $(\varepsilon, t_1, \ldots, t_n)$ such that $E = V(t_1, \ldots, t_r)$ and $f^{\#}(\pi) = \varepsilon t_1^{d_1}, \ldots, t_r^{d_r}$. In particular, $X \to B$ locally factors through a morphism $h: X \to X_0 = \operatorname{Spec}(k[\varepsilon, t_1, \ldots, t_n]/(\varepsilon^n))$, and if $X \to B$ is of finite type, then h is étale. 3.3.2. The main result. Now we can prove our main result on resolution of *B*-schemes.

Theorem 3.3.3. Let k be a field of characteristic zero and $B = \operatorname{Spec}(k[\pi]/(\pi^n))$. Then there exists a procedure \mathcal{F} which obtains as an input a generically smooth morphism of finite type $X \to B$ and a closed nowhere dense subscheme $Z \to X$, and outputs a modification $\mathcal{F}(X, Z) \colon X' \to X$ with a X' a ptm and a boundary E' on X' such that $Z' = Z \times_X X'$ is an E'-monomial divisor and (X', E') is a distinguished B-pair. In addition, \mathcal{F} depends smooth functorially on the input: if $Y \to X$ is smooth and $T = Z \times_X Y$, then $\mathcal{F}(Y, T) = \mathcal{F}(X, Z) \times_X Y$.

Proof. First, by Theorem 2.3.1 there exists a modification $X_1 \to X$ with X_1 a ptm. Second, applying Theorem 3.1.9 to X_1 and $Z_1 \times_X X_1$ we obtain a modification of ptm's $X_2 \to X_1$ and an ordered boundary E_2 on X_2 such that $Z_2 = Z \times_X X_2$ is E_2 -monomial. Third, the generic smoothness of f implies that generically π generates the nilradical $\mathcal{N}_2 \subset \mathcal{O}_{X_2}$ on an open subscheme U. Principalizing $X_2 \setminus U$ on (X_2, E_2) by Theorem 3.1.9 we obtain a modification $X_3 \to X_2$, such that X_3 is provided with an ordered boundary E_3 and $\mathcal{N}_3|_U = (\pi)|_U$ for $U = X_3 \setminus E_3$.

Now, locally on X_3 there exists a factorization $(\pi) = \mathcal{N}_3 \mathcal{I}$, and the locally principal ideal \mathcal{I} is invertible because it is trivial at the generic point and X_3 has no embedded components. Note that \mathcal{I} does not have to be E_3 -monomial, moreover it is not uniquely defined hence the factorization does not have to globalize. However, blowing up the components of E_3 as in Lemma 3.2.4 we obtain a modification $X' \to X_3$ with a trivial reduction such that the pullbacks of each \mathcal{I} as above become monomial, that is, on X' we locally have factorizations $(\pi) = \mathcal{N}'\mathcal{I}$ with an E'-monomial \mathcal{I} . Such a factorization is unique because if $(\pi) = \mathcal{N}'\mathcal{I}'$ is another factorization with an E_3 -monomial \mathcal{I}' , then $V(\mathcal{I})$ and $V(\mathcal{I}')$ restrict to the same divisor on \widetilde{X}_3 and hence contain the same components of E_3 with the same multiplicities. Thus the factorization $(\pi) = \mathcal{N}\mathcal{I}$ globalizes, and this shows that $X' \to X$ is as required.

The smooth-functoriality of the construction follows from functoriality of all ingredients we used. $\hfill \Box$

3.4. **Other geometric spaces.** We deduced our results from the classical principalization theorem for algebraic varieties. Now, let us discuss what can be done for more general schemes, complex analytic spaces, etc. These are nowadays standard reductions that were described in details in a few papers, so we only outline them.

3.4.1. Analytic spaces. It is folklore knowledge that the classical principalization theorem and its proof apply almost verbatim for complex analytic spaces and non-archimedean analytic spaces of characteristic zero, though, to the best of our knowledge, the only formalized framework, in which this was checked, is given in [BM97]. Starting with principalization for analytic spaces one can establish analogues of Theorems 2.3.1, 3.1.9 and 4.4.1 using precisely the same arguments as in the paper.

3.4.2. Schemes with enough derivations. In fact, the classical principalization algorithm and its justification applies to the wider context of ideals on regular schemes with enough derivations, where the latter were defined in [Tem12, Remark 1.3.2(iii)]. The only new (and simple) thing one should check is that any smooth blow up of an X with enough derivations also has enough derivations. Strictly speaking, this was also never published, but a much more involved case of an algorithm in the

setting of enough derivations was worked out with all details in [ATW20]. Moreover, the functoriality holds with respect to arbitrary regular morphisms and not only the smooth ones. Once principalization for such schemes is known, our arguments extend to this setting verbatim, and one obtains that Theorems 2.3.1, 3.1.9 and 4.4.1 hold for schemes locally embeddable into regular schemes with enough derivations, and the functoriality is with respect to arbitrary regular morphisms. Then the usual formal argument (see [Tem12, §5.2]) implies that the same results extend to analytic spaces (recovering the claim of §3.4.1) and formal schemes locally embeddable into regular formal schemes with enough derivations.

3.4.3. Arbitrary qe schemes. It seems certain that the maximal generality, in which Theorems 2.3.1, 3.1.9 and 4.4.1 should hold, is that of arbitrary qe schemes of characteristic zero. However, resolution of non-reduced schemes was not established in such generality and only a weaker form of principalization was proved in [Tem18, Theorem 1.1.11] – it outputs a regular scheme with a monomial ideal, but the intermediate steps can be singular. This would suffice to deduce the last two theorems, but it cannot be used for the non-embedded resolution of generically nil-principal schemes. All in all, the only stumbling block is resolution of arbitrary non-reduced qe schemes.

4. Complements

In this section we will slightly refine our main result by constructing a global embedding of a modification into a log smooth *B*-scheme. This requires an additional preparation.

4.1. **Factorization.** In addition to the analogues of the usual resolution and principalization we will need certain results about factorization and resolution of indeterminacies for ptm's. We will only cover the cases that will be needed in our applications.

Theorem 4.1.1. Assume that $f: Y \to X$ is a modification of ptm's with a trivial reduction. Assume, in addition, that there exist an ordered boundary $E = \sum_{i=1}^{l} E_i$ on X such that f is an isomorphism over $X \setminus E$. Then there exists a sequence of blowings up of components of \tilde{E}_i 's $Y' \to Y$ such that the composed morphism $f': Y' \to X$ splits into a composition of blowings up of components of \tilde{E}_i 's. The sequence $Y' \to Y$ and the splitting of f' depend on the datum functorially with respect to smooth morphisms $\overline{X} \to X$ and pullbacks $\overline{f}: \overline{Y} = Y \times_X \overline{X} \to \overline{X}$ and $\overline{E} = E \times_X \overline{X}$.

Proof. To simplify notation we identify the underlying topological spaces of X and Y. Notice that locally we can factor $\mathcal{N}_X \mathcal{O}_Y$ as $\mathcal{I} \mathcal{N}_Y$, where \mathcal{I} is invertible – just take local generators ε_Y and ε_X , then $\varepsilon_X = a\varepsilon_Y$ and we take $\mathcal{I} = (a)$. This does not yield a global divisor $\mathcal{N}_Y/\mathcal{N}_X$ because an element a is defined only up to adding an element annihilated by ε_Y , but the reduction (\tilde{a}) is unique and we obtain a global divisor on \tilde{X} that we denote $\tilde{D} = \mathcal{N}_Y/\mathcal{N}_X$. Clearly, it is \tilde{E} -monomoial, so we fix the decomposition $\tilde{D} = \sum_{i,n} n \tilde{D}_{i,n}$ with $\tilde{D}_{i,n} \subset \tilde{E}_i$. Now, let $Y' \dashrightarrow Y$ be the sequence obtained by successive blowing up the centers $V_{i,j} = \coprod_{n \ge j} \tilde{D}_{i,n}$ and let E' be the transform of the boundary to Y'. By Lemma 3.2.4 the local divisors (a) become E'-monomial on Y'. The reductions (\tilde{a}) agree on the intersections (this happens

already on X), hence we obtain a factorization $\mathcal{N}_{Y'} = D' \mathcal{N}_X$ with an E'-monomial divisor D'.

Choose the presentation $D' = \sum_{i,n} nD'_{i,n}$ with $D'_{i,n} \subset E_i$ (we do not need this, but in fact only even n's show up and $\widetilde{D}'_{i,2n} = \widetilde{D}_{i,n}$). We claim that $Y' \to X$ splits into the blowing up sequence $Y' = X_m \dashrightarrow X_0 = X$ with centers $V'_{i,j} = \prod_{n\geq j} \widetilde{D}'_{i,n}$. Indeed, this can be checked locally at a point $x \in X$. Choose parameters $\varepsilon_X, t_1, \ldots, t_n \in \mathcal{O}_{X,x}$ such that $E = V(t_1 \ldots t_r)$ and let n_i be such that $x \in D'_{i,n_i}$. Then $\varepsilon_Y = \varepsilon_X / \prod_{i=1}^r t_i^{n_i}$ is a generator of the nilradical of $\mathcal{O}_{Y',x}$ and hence an embedding $\mathcal{O}_{X_m,x} \subseteq \mathcal{O}_{Y',x}$ arises. Moreover, it is an isomorphism by Nakayama's lemma because the reductions coincide and $\mathcal{O}_{X_m,x}$ contains a generator of the nilradical of $\mathcal{O}_{y',x}$. Thus, $Y' = X_m$ and we have constructed a splitting of $Y' \to X$. The check that our construction is smooth-functorial is straightforward and we omit it. \Box

Corollary 4.1.2. Let $X \to B_n = \operatorname{Spec}(k[\pi]/(\pi^n))$ be generically smooth. Then there exists a construction $\mathcal{G}(X \to B_n) = Y \to X \to B_n$ such that the composed morphism is smooth away from an snc divisor. This construction is functorial with respect to smooth morphism.

Proof. As $X \to B_n$ is generically smooth, X is generically nilprincipal and we may apply Theorem 2.3.1 to get $X' \to X$ such that X is a ptm. Let $Z \subset X$ be the singular locus of the composed morphism $X' \to X \to B_n$, and by Theorem 3.1.9 we have a principalization $\mathcal{P}(X', \emptyset, Z) : Y \to X'$ such that the composed morphism $Y \to B_n$ is smooth away from an snc divisor. This process is smooth-functorial as both ingredients are.

Remark 4.1.3. We make two remarks on this result:

- (1) The morphism $Y \to X'$ is a modification of ptms with trivial reduction, so if we wish to apply Theorem 4.1.1 we may.
- (2) In practice, this means that given a generically smooth $X \to B_n$ there is a resolution which étale-locally looks like $X' = \operatorname{Speck}[\epsilon, t_1, ..., t_l]/(\epsilon^n) \to B_n$ given by $\pi \mapsto \epsilon \cdot t_1^{n_1} ... t_l^{a_l}$.

4.2. **Retracts.** By a retract of a ptm X we mean a morphism $r: X \to \widetilde{X}$ which is a retract of the closed immersion $i: \widetilde{X} \hookrightarrow X$, that is, $r \circ i = \operatorname{id}_{\widetilde{X}}$. In general, a retract does not exist, but if X is of finite type over a field k, then it exists whenever X is affine. Indeed, X is a nilpotent thickening of \widetilde{X} and \widetilde{X} is smooth over k, hence the k-morphism $\operatorname{id}_{\widetilde{X}}$ lifts to a k-morphism $r: X \to \widetilde{X}$. (For an arbitrary ptm X, a retract at least exists at the generic points $\eta \in X$, namely, $r_{\eta}: \eta \to \widetilde{\eta}$ corresponds to a choice of a field of definition $k(\widetilde{\eta}) \hookrightarrow \mathcal{O}_{X,\eta}$ of the Artin ring $\mathcal{O}_{X,\eta}$.) We will show that after blowing up X one can extend a generic retract to a retract of the whole X. The argument will be similar to the proof of Theorem 4.1.1 but a bit more elaborate. Here is the main particular case we will need.

Lemma 4.2.1. Let X be a ptm of finite type over a field k with nilradical $\mathcal{N} \subset \mathcal{O}_X$, let $D = V(\mathcal{I})$ be a Cartier divisor with a smooth reduction $\widetilde{D} = D \times_X \widetilde{X}$ and complement $X_0 = X \setminus D$, and let $i: X_0 \hookrightarrow X$ denote the open immersion. Assume that D is irreducible with generic point ν and $r_0: X_0 \to \widetilde{X}_0$ is a retract of X_0 , then

(i) The retract r_0 extends to r if and only if the homomorphism $\phi: i_*\mathcal{O}_{\widetilde{X}_0} \hookrightarrow i_*\mathcal{O}_{X_0}$ induced by r_0 takes $\mathcal{O}_{\widetilde{X}}$ to \mathcal{O}_X .

(ii) The following three conditions are equivalent and they hold for a large enough n: (a) ϕ takes $\mathcal{O}_{\widetilde{X}}$ to $\mathcal{O}_X[\mathcal{I}^{-n}\mathcal{N}]$, (b) ϕ takes $\mathcal{O}_{\widetilde{X},\nu}$ to $\mathcal{O}_{X,\nu}[\mathcal{I}_{\nu}^{-n}\mathcal{N}_{\nu}]$, (c) r_0 extends to a global retract on $\operatorname{Bl}_{\widetilde{nD}}(X)$.

(iii) Assume that $n \geq 1$ is such that the conditions in (ii) are satisfied and let $X' = \operatorname{Bl}_{\widetilde{D}}(X)$. Then ϕ takes $\mathcal{O}_{\widetilde{X}'}$ to $\mathcal{I}^{1-n}\mathcal{O}_{X'}$.

Proof. All claims can be checked locally on X, where we use in (ii) that X is quasicompact. Thus, we can assume that X = Spec(A) and $\mathcal{I} = (t)$, and then $X_0 =$ $\text{Spec}(A_t)$, $\tilde{X} = \text{Spec}(\tilde{A})$ and $\tilde{X}_0 = \text{Spec}(\tilde{A}_t)$, where $\tilde{A} = A/\mathcal{N}$. Let $s_0 \colon \tilde{A}_t \to A_t$ be the section corresponding to r_0 . Since $\tilde{X} \to X$ is a homeomorphism, r_0 extends to r if and only if the homomorphism of sheaves extends, that is $s_0(\tilde{A}) \subseteq A$. This proves (i).

Recall that $X' = \operatorname{Bl}_{n\overline{D}}(X) = \operatorname{Bl}_{(t^n)+\mathcal{N}}(X)$ consists of the *t*-chart, hence $X' = \operatorname{Spec}(A_n)$, where $A_n = A[t^{-n}\mathcal{N}]$. Clearly, $X' \to X$ is a homeomorphism and we denote the preimage of ν by ν' . Applying (i) to X' we obtain that (a) and (c) are equivalent, and the equivalence of (a) and (b) reduces to the claim that if r_0 extends to ν' , then r_0 extends to the whole X'. But X' can be obtained from the ptm X by n successive blowings up of \widetilde{D} , hence X' is a ptm and any regular function on X_0 that extends to ν' automatically extends to the whole X' by Corollary 2.2.7. The claim about extension of r_0 follows.

Let us now prove that for a large n the condition (a) is satisfied. Given $\tilde{a} \in \tilde{A}$ choose a lift $a \in A$ and notice that $s_0(\tilde{a}) - a \in \mathcal{N}_t$ and hence lies in a subring $A_n = A[t^{-n}\mathcal{N}]$ of A_t . Choosing a set $\tilde{a}_1, \ldots, \tilde{a}_l$ of k-generators of \tilde{A} and a large enough n, we obtain that $s_0(\tilde{a}_i) \in A_n$ for any $1 \leq i \leq l$ and hence $s_0(\tilde{A}) \subset A_n$. This proves the second part of (ii).

Finally, in (iii) X' = Spec(A'), where $A' = A_1$, and hence the nilradical of A' is $\mathcal{N}' = t^{-1}\mathcal{N}$ and $\widetilde{A} = \widetilde{A}'$. Therefore, $s_0(\widetilde{A}') \subset A_n = A'[t^{1-n}\mathcal{N}']$, proving (iii). \Box

Theorem 4.2.2. There exists a method which to any ptm with a boundary (X, E) of finite type over a field k of characteristic zero and a generic retract $r_{\eta} : \eta \to \tilde{\eta}$, where $\eta = \coprod_{i=i}^{n} \operatorname{Spec}(\mathcal{O}_{X,\eta_i})$ and η_1, \ldots, η_n is the set of generic points of X, associates a sequence of admissible blowings up

$$\mathcal{R}(X, E, r_{\eta}) \colon (X', E') \dashrightarrow (X, E)$$

such that r_{η} extends to a retract $r: X' \to \widetilde{X}'$. Moreover, one can construct \mathcal{R} functorially with respect to étale morphisms $Y \to X$ provided with a compatible pair of generic retracts r_{η_X} and r_{η_X} .

Proof. We will first construct the retract as a composition of two admissible blowings up sequences, and only then discuss the functoriality. To simplify notation we assume that X is irreducible and η is the generic point, but the argument works in general as well. Let $s_{\eta} \colon k(\eta) \hookrightarrow \mathcal{O}_{X,\eta}$ be the field of coefficients corresponding to r_{η} . Recall that X has no embedded components, hence $\mathcal{O}_{X,x} \subseteq \mathcal{O}_{X,\eta}$ for any $x \in X$ and by Lemma 4.2.1(i) r_{η} extends to r locally at $x \in X$ if and only if $s_{\eta}(\mathcal{O}_{\tilde{X},x}) \subseteq \mathcal{O}_{X,x}$. Clearly, this is an open condition and we denote by U the open set of such points $x \in X$, and denote by $Z = X \setminus U$ the complement. In other words, Z is the indeterminacy locus of the rational map $r_{\eta} \colon X \to \tilde{X}$. The first part of \mathcal{R} is the principalization $\mathcal{P}(X, E, Z) \colon X' = X_l \dashrightarrow X_0 = X$ of Z. Let E'_1, \ldots, E'_l be the components of E' created by \mathcal{P} , then the indeterminacy locus Z' of $r_{\eta} \colon X' \to \widetilde{X}'$ is contained in $E'' = \sum_{i=1}^{l} E'_{i}$.

The remaining procedure iteratively blows up unions of components of \widetilde{E}'' . In particular, the reduction \widetilde{X}' remains unchanged, so we can identify \widetilde{E}'_i with subschemes of further blowings up. The choice of the center on each step is done by the following rule: for any generic point $\nu \in E''$ let $i = i(\nu)$ be the number for which $\nu \in E'_i$ and let $n(\nu)$ be the minimal $n \geq 0$ which satisfies the condition of Lemma 4.2.1(ii). The step takes all generic points ν such that $n(\nu) > 0$ (i.e. the point is in the indeterminacy locus) and the lexicographic value of $(n(\nu), i(\nu))$ is maximal, and blows up the reduced closed subschemes with these generic points. In other words, we blow up components of \widetilde{E}'_i with the maximal value of the invariant $(n(\nu), i(\nu))$. The maximal value of the invariant drops after each blowing up by Lemma 4.2.1(ii) and (iii). Part (i) of the same lemma says that once n drops to zero, the indeterminacy locus is empty, that is, r_η extends to a global retract.

Finally, let us discuss the functoriality. So let us assume that $f: Y \to X$ is étale, $E_Y = E \times_X Y$ and a generic retract r_{η_Y} is the pullback of r_{η_X} . The principalization is functorial with respect to smooth morphisms, so we should only check that the indeterminacy locus and the invariant $n(\nu)$ used in the second step are compatible with f. Moreover, if D is the irreducible component of ν , then by Lemma 4.2.1(ii) $n(\nu)$ is the minimal n such that blowing up \overline{nD} resolves the indeterminacy at ν . Blowings up are compatible with smooth (even flat) morphisms, hence it suffices to check functoriality of the indeterminacy locus: if $y \in Y$ and x = f(y), then r_{η_X} extends to x if and only if r_{η_Y} extends to y. Furthermore, since the indeterminacy locus is open, localizing X and Y appropriately we can reduce the problem to the following particular case: $r_{X_{\eta}}$ extends to a global retract r_X if and only if $r_{\gamma_{\eta}}$ extends to a global retract r_Y . Clearly, it suffices to check the latter locally, so we can assume that X and Y are affine.

Assume that $r_{\eta_X} : \eta_X \to \eta_{\widetilde{X}}$ extends to $r : X \to \widetilde{X}$. Then a morphism $Y \to X \to \widetilde{X}$ arises and the \widetilde{X} -morphism $\operatorname{id}_{\widetilde{Y}}$ lifts uniquely to a \widetilde{X} -morphism $r_Y : Y \to \widetilde{Y}$ by the étalenes of \widetilde{Y} over \widetilde{X} and affineness of Y. Thus, the retract r_X lifts uniquely to a retract r_Y . Since the lift r_{η_Y} of r_{η_X} is also unique by the same argument, both agree, that is r_Y is the extension of r_{η_Y} . Conversely, if r_{η_Y} extends to r_Y , then it is easy to see that r_{η_X} extends to r_X .

We do not know if functoriality holds for smooth morphisms as well.

4.3. Log blowings ups. We will need one more construction on a ptm with boundary. Informally speaking, it replaces boundary components by varieties intersecting X so that one obtains a thick version of a normal crossings variety. The most convenient way to realize this is by using log geometry.

4.3.1. Associated log structure. As often happens one can encode the boundary in a log structure, and in various applications this is even a more conceptual approach. Given a ptm with a boundary (X, E) we define the associated log structure M_X as follows: for a point $x \in X$ let $t_0 = \varepsilon$ be a generator of the nilradical of \mathcal{O}_x and let $t_1, \ldots, t_r \in \mathcal{O}_x$ be parameters that define the irreducible components of E passing through x and consider the prelogarithmic structure $\overline{M}_x = \bigoplus_{i=0}^r \mathbb{N} \log(t_i) \to \mathcal{O}_x$ which sends $\log(t_i)$ to t_i . Since the choice of ε and t_i is unique up to a unit, the

associated log structure M_x is independent of the choices, and it is easy to see that the local definitions patch to a global log structure M_X on X. Clearly, E can be reconstructed from M_X , so we also address to the log scheme (X, M_X) as a *ptm* with a boundary.

4.3.2. Log blowings up of reduced divisors. We will use a very special form of log blowings up of a ptm with a boundary (X, M_X) .

Lemma 4.3.3. Let (X, M_X) be a ptm with a boundary and D a smooth component of the induced boundary \tilde{E} of \tilde{X} . Then $X'' = \operatorname{Bl}_{\tilde{D}}(X)$ is an irreducible component of $X' = \operatorname{LogBl}_{\tilde{D}}(X)$. In addition, X'' is a monoidal subscheme of X'.

Proof. The claim can be checked locally over a point $x \in \widetilde{D}$, so assume that $X = \operatorname{Spec}(A)$ and $(\varepsilon = t_0, t = t_1, \ldots, t_n)$ is a family of parameters such that $\widetilde{D} = (\varepsilon, t)$ and $E = V(t_1 \ldots t_r)$ at x. Then $Y = \operatorname{Spec}(\mathbb{Z}[t_0, \ldots, t_r])$ is a chart for X, and hence $X' \to X$ is the base change of the log blowing up of Y along (t_0, t_1) . Since Y is log regular, the latter is the usual blowing up described by the charts, and hence X' is glued from two charts: $X'_t = \operatorname{Spec}(A[\varepsilon']/(t\varepsilon' - \varepsilon))$ and $X'_{\varepsilon} = \operatorname{Spec}(A[t']/(t'\varepsilon - t))$.

Recall that $X'' = \operatorname{Spec}(A'')$, where $A'' = A[\frac{\varepsilon}{t}]$, and note that A'' is the image of $A' = A[\varepsilon']/(t\varepsilon' - \varepsilon)$ in A_t , hence X'' is the irreducible component of X', is the strict transform of X and lies in the *t*-chart. Furthermore, it is easy to see that $A'' = A'/(\varepsilon'^h)$, where h is the thickness of X, hence X'' is given by vanishing of the monomial ideal (ε'^h) .

Here is a concrete typical example.

Example 4.3.4. Let $A = k[t, \varepsilon]/(\varepsilon^n)$ and X = Spec(A) with the log structure induced by $\mathbb{N}\log(\varepsilon) \oplus \mathbb{N}\log(t)$. Then $\widetilde{D} = V(t, \varepsilon)$ is a monomial subscheme, which is also a smooth divisor in \widetilde{X} . The log blowing up $X' = \text{LogBl}_{\widetilde{D}}(X)$ is glued from two charts:

$$X'_t = \operatorname{Spec}(A[\varepsilon']/(t\varepsilon' - \varepsilon)) = \operatorname{Spec}(k[\varepsilon', t]/(t^2\varepsilon'^2))$$

and

$$X'_{\varepsilon} = \operatorname{Spec}(A[t']/(t'\varepsilon - t)) = \operatorname{Spec}(k[t', \varepsilon]/\varepsilon^2)$$

Thus, X' has two irreducible components. The first component is $\operatorname{Bl}_{\widetilde{D}}(X) = \operatorname{Spec}(k[t,\varepsilon'])/(\varepsilon'^2)$; it is contained in X'_t . The other component is a ptm with nilradical (ε) and reduction isomorphic to \mathbf{P}^1_k with coordinate t'; it is contracted to \widetilde{D} in X.

For our applications we will need to construct sequences of log blowings. The following results follows from Lemma 4.3.3 by an obvious induction on the length of the sequence.

Corollary 4.3.5. Let (X, M_X) be a log ptm with a boundary, and let D_1, \ldots, D_l be smooth components of the associated snc boundary $\tilde{E} \subset \tilde{X}$. Then there exists a sequence of log blowings up $X_l \dashrightarrow X_0 = (X, M_X)$ such that the following conditions are satisfied: the strict transform X'_i of X in X_i is a monomial subscheme of X_i , each X'_{i+1} is the blowing up of X'_i along \tilde{D}_i , in particular, each $X'_i \to X$ is a modification with the trivial reduction and $\tilde{D}_i \hookrightarrow X'_i \to X_i$ are monomial subschemes, and $X_{i+1} = \text{LogBl}_{\tilde{D}_i}(X_i)$. 4.4. Log smooth resolution. To get our last main result it remains to assembly various pieces we have developed altogether.

Theorem 4.4.1. Assume that k is a field of characteristic zero, $B = \operatorname{Spec}(k[\pi]/(\pi^n))$ and $X \to B$ is a generically smooth morphism of finite type with a closed nowhere dense subscheme $Z \hookrightarrow X$. Let M_B be the log structure on B generated by $\mathbb{N}\log(\pi) \to M_B$, where $\log(\pi)$ is mapped to π . Then there exists a log smooth morphism $(Y, M_Y) \to (B, M_B)$, an irreducible component X' of Y and modification $X' \to X$, such that $Z' = Z \times_X X'$ is a divisor and both X' and Z' are monomial subschemes of Y. In addition, the choice of Y and X' can be made depending étale-functorially only on $X \to B$, $Z \hookrightarrow X$ and a generic retract $r_{\eta} \colon \eta_X \to \eta_{\widetilde{X}}$, and if $X \to B$ is proper one can also choose $Y \to B$ to be proper.

Proof. Acting as in the proof of Theorem 3.3.3 we can find a modification $X_1 \to X$ such that X_1 is a ptm with an ordered boundary E_1 and $Z_1 = Z \times_X X_1$ is E_1 monomial. Next, the generic smoothness of f implies that generically π generates the nilradical $\mathcal{N} \subset \mathcal{O}_{X_1}$ on an open subscheme U. Principalizing $X_2 \setminus U$ on (X_2, E_2) by Theorem 3.1.9 we obtain a modification $X_2 \to X_1$, such that $\mathcal{N}|_U = (\pi)|_U$ for $U = X_2 \setminus E_2$. Next, we apply Theorem 4.2.2 to construct a modification $g: X' \to X_2$ with a ptm X' and an ordered boundary E' containing the preimage of E_2 such that X' possesses a retract $X' \to \tilde{X}'$.

Set $U' = X' \setminus E'$ and $Y' = \tilde{X}' \times_{\tilde{B}} B = \operatorname{Spec}_{\tilde{X}'} \mathcal{O}'_{\tilde{X}}[\pi]/(\pi^n)$. Then the natural morphism $f: X' \to Y'$ is a modification with a trivial reduction. In addition, f is an isomorphism over the complement to \tilde{E}' , hence by Theorem 4.1.1 replacing X' once again by a successive blowing up along components of E', we can achieve that f factors into a composition $X' = Y'_n \dashrightarrow Y'_0 = Y'$ of blowings up of components of \tilde{E}'_i 's, say $Y'_{i+1} = \operatorname{Bl}_{V_i}(Y'_i)$.

Finally, we provide Y' with the log structure combined from the log structures induced from B and the log structure on $\widetilde{Y}' = \widetilde{X}'$ given by \widetilde{E}' . More formally, consider the closed immersion $i: \widetilde{U}' \hookrightarrow \widetilde{Y}'$, then $M_{\widetilde{Y}'} = i_* \mathcal{O}_{\widetilde{U}'}^{\times} \cap \mathcal{O}_{\widetilde{Y}'}$ provides \widetilde{X} with the structure of a log smooth log scheme over k (with the trivial log structure $k^{\times} \hookrightarrow$ k) and we set $(Y', M_{Y'}) = (\widetilde{Y}', M_{\widetilde{Y}'}) \times_{\operatorname{Spec}(k)} (B, M_B)$. In particular, the projection $(Y, M_{Y'}) \to (B, M_B)$ is log smooth. By Corollary 4.3.5 there exists a sequence of log blowings up $Y = Y_n \dashrightarrow Y_0 = Y'$ such that $Y'_i \hookrightarrow Y_i$ is the irreducible component which is a modification of Y' and the centers are log subschemes $V_i \hookrightarrow \widetilde{Y}' = \widetilde{Y}'_i \hookrightarrow$ $Y'_i \hookrightarrow Y_i$. Thus, X' is an irreducible component of Y and since the log blowings up are log étale the composition $Y \dashrightarrow Y' \to B$ is log smooth.

The étale-functoriality of the construction follows from functoriality of all ingredients we used. $\hfill \Box$

References

- [AM69] M. F. Atiyah and I. G. Macdonald, Introduction to commutative algebra, Addison-Wesley Publishing Co., Reading, Mass.-London-Don Mills, Ont., 1969. MR 0242802 (39 #4129)
- [ATW20] D. Abramovich, M. Temkin, and J. Włodarczyk, Relative desingularisation and principalization of ideals, March 2020, http://arxiv.org/abs/2003.03659.
- [BM97] Edward Bierstone and Pierre D. Milman, Canonical desingularization in characteristic zero by blowing up the maximum strata of a local invariant, Invent. Math. 128 (1997), no. 2, 207–302. MR 1440306 (98e:14010)

- [BM08] _____, Functoriality in resolution of singularities, Publ. Res. Inst. Math. Sci. 44 (2008), no. 2, 609–639. MR 2426359
- [BMT11] Edward Bierstone, Pierre D. Milman, and Michael Temkin, Q-universal desingularization, Asian J. Math. 15 (2011), no. 2, 229–249. MR 2838221
- [Gro67] Alexander Grothendieck, Éléments de géométrie algébrique., Inst. Hautes Études Sci. Publ. Math. (1960-1967).
- [Hir64] Heisuke Hironaka, Resolution of singularities of an algebraic variety over a field of characteristic zero. I, II, Ann. of Math. (2) 79 (1964), 109–203; ibid. (2) 79 (1964), 205–326. MR 0199184
- [HK04] Brendan Hassett and Sándor J. Kovács, Reflexive pull-backs and base extension, J. Algebraic Geom. 13 (2004), no. 2, 233–247. MR 2047697
- [Kol07] János Kollár, Lectures on resolution of singularities, Annals of Mathematics Studies, vol. 166, Princeton University Press, Princeton, NJ, 2007. MR 2289519 (2008f:14026)
- [Tem12] Michael Temkin, Functorial desingularization of quasi-excellent schemes in characteristic zero: the nonembedded case, Duke Math. J. 161 (2012), no. 11, 2207–2254. MR 2957701
- [Tem18] _____, Functorial desingularization over q: boundaries and the embedded case, Israel J. Math. 224 (2018), no. 1, 455–504. MR 3799764
- [Wło05] Jarosław Włodarczyk, Simple Hironaka resolution in characteristic zero, J. Amer. Math. Soc. 18 (2005), no. 4, 779–822 (electronic). MR 2163383

EINSTEIN INSTITUTE OF MATHEMATICS, THE HEBREW UNIVERSITY OF JERUSALEM, EDMOND J. SAFRA CAMPUS, GIV'AT RAM, JERUSALEM, 91904, ISRAEL

Email address: alex.motzkin@mail.huji.ac.il Email address: michael.temkin@mail.huji.ac.il